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JC 598 U.S. PTO

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JC 525 U.S. PTO
09/491389
01/26/00

Docket No. : ARC-P002

APPLICATION TRANSMITTAL LETTER

Assistant Commissioner of Patents
United States Patent and Trademark Office
Washington, D.C. 20231
ATTN: BOX PATENT APPLICATION

Sir:

Transmitted herewith for filing is the patent application of

Inventor(s): Ku, et al.
Entitled: Method and Apparatus of Relative Datapath Cell Placement with Structure Bonding and Exponential Field Sampling

26 No. pages of specification, including title page, claims and abstract
9 No. sheets of X informal, formal drawings

Also enclosed are:

- Executed Combined Declaration and Power of Attorney for Patent Application
 An Original Executed Assignment of the Application
 Form PTO-1595 (Recordation Cover Sheet for Assignment)
 Verified Statement Claiming Small Entity Status with Cover Sheet
 An Information Disclosure Statement (Form PTO-1449A and Form PTO-1449B)
 A copy of References cited in Information Disclosure: 4 documents

FEES DUE

The fees due for filing the application pursuant to 37 C.F.R. 1.16 and for recording the Assignment, if any, are determined as follow:

CLAIMS					
	No. of Claims		Extra Claims	Rate	Fees
Basic Application Fee (\$690 large entity; \$345 small entity)					\$ 345.00
Total Claims	24	Minus 20 =	4	X \$18 = X \$ 9 (small) =	36.00
Total Independent Claims	5	Minus 3 =	2	X \$78 = X \$39 (small) =	78.00
If Multiple Dependent Claims are presented, add \$260.00 or \$130.00 (small)					
If Assignment enclosed, add Assignment Recording Fee \$40.00					40.00
TOTAL APPLICATION FEE DUE					\$ 499.00

PAYMENT OF FEES

The full fee due in connection with this communication is \$ 499.00
and is provided as follows:

- The Commissioner is hereby authorized to charge the fees associated with this communication or credit any overpayment to **Deposit Account No: 500482**. A duplicate copy of this authorization is enclosed.
- A Check No. 6022 for the above specified full fee is enclosed. However, in case Applicant inadvertently miscalculated any required fee, the Commissioner is hereby authorized to charge the necessary additional amount associated with this communication or credit any overpayment to **Deposit Account No: 500482**. A duplicate copy of this authorization is enclosed.

This application is filed pursuant to 37 C.F.R. 1.53 in the name of the above-identified Inventor(s).

Please direct all correspondence concerning the above-identified application to the following address:

FERNANDEZ & ASSOCIATES LLP
PATENT ATTORNEYS
PO BOX D
MENLO PARK, CA 94026-6204

(650) 325-4999
(650) 325-1203 : FAX
EMAIL: *iploft@iploft.com*

Respectfully submitted,


PETER C. SU, ESQ.

Reg. No. 43,939

1/26/2000
Date

The Assistant Commissioner of Patents
United States Patent and Trademark Office
Washington, D.C. 20231
ATTN: Box Patent Application

Re: U.S. Utility Patent Application
Appl. No. (Not yet assigned); Filed 1/26/2000
For: **Method and Apparatus of Relative Datapath Cell Placement with Structure
Bonding and Exponential Field Sampling**
Inventor(s): **Ku, et al.**
Docket No.: ARC-P002

Sir:

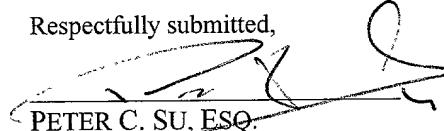
The following documents are forwarded herewith for action by the U.S. Patent and Trademark Office:

1. U.S. UTILITY APPLICATION
entitled: **Method and Apparatus of Relative Datapath Cell Placement with Structure
Bonding and Exponential Field Sampling**
having named inventor(s):
Ku, et al.
 - a. a specification consisting:
 - (i) 20 pages prior to the claims, including title page;
 - (ii) 5 pages of claims;
 - (iii) 1 page abstract;
 - b. 9 sheets of informal drawings: (FIGs.);
2. An original, executed Combined Declaration and Power of Attorney by named inventors;
3. Form PTO-1082 (in duplicate);
4. Cover letter for Assignment (Form PTO-1595)
5. An original, executed Assignment to Arcadia Design Systems., executed by named inventors, recordation of which is hereby requested;
6. A return post card; and
7. Check No. 6022 for \$ 499 to cover:

Patent application filing fee:	\$ 459.00
Assignment Recordation fee:	\$ 40.00
Excess claims fee:	\$.00
8. Verified Small Entity Status Statement with Cover Sheet
9. Information Disclosure Statement and Form PTO 1449

It is respectfully requested that the attached postcard be stamped with the filing date of the above documents and unofficial application number and returned to the addressee as soon as possible.

1/26/2000
Date

Respectfully submitted,

PETER C. SU, ESQ.
Reg. No. 43,939

FERNANDEZ & ASSOCIATES LLP
PATENT ATTORNEYS
PO BOX D
MENLO PARK, CA 94026-6204

(650) 325-4999
(650) 325-1203 : FAX
EMAIL: *iploft@iploft.com*

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

In re application of :

Application No.: (Not yet assigned)
Filed: 1/26/ 2000

Group No.: (Not yet assigned)
Examiner: (Not yet assigned)

Title: **Method and Apparatus of Relative Datapath Cell Placement with Structure Bonding and Exponential Field Sampling**

Inventor(s): **Ku, T. et al.**

Assistant Commissioner for Patents
Washington, D.C. 20231

**SUBMISSION OF VERIFIED STATEMENT
TO ESTABLISH SMALL ENTITY STATUS**

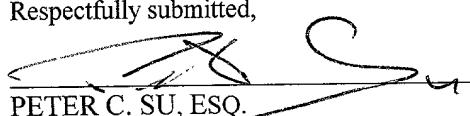
The attached statement is being submitted to establish small entity status in this

Application
 Patent

by the

- Independent inventor(s) 37 CFR 1.9(c) and 1.27(b)
 Non-inventor supporting claim by another 37 CFR 1.9(c) and 1.27(b)
 Small Business Concern 37 CFR 1.9(d) and 1.27(c)
 Nonprofit Organization 37 CFR 1.9(e) and 1.27(d)

Respectfully submitted,


PETER C. SU, ESQ.
Reg. No. 43,939

FERNANDEZ & ASSOCIATES LLP
PATENT ATTORNEYS
PO BOX D
MENLO PARK, CA 94026-6204

(650) 325-4999
(650) 325-1203 : FAX
EMAIL: *iploft@iploft.com*

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

Application No.: (Not yet assigned)
Filed on: 1/26/2000Patent No.: (Not yet assigned)
Issued on: (Not yet assigned)Title: **Method and Apparatus of Relative Datapath Cell Placement with Structure Bonding and Exponential Field Sampling**
Inventor(s): **Ku, et al.****VERIFIED STATEMENT CLAIMING SMALL ENTITY STATUS**
(37 CFR 1.9(f) and 1.27(c) -- SMALL BUSINESS CONCERN)

I hereby declare that I am

- the owner of the small business concern identified below:
- an official of the small business concern empowered to act on behalf of the concern identified below:

Name of Small Concern:
ARCADIA DESIGN SYSTEMSAddress of Small Concern:
**2801 Orchard Parkway
San Jose, CA 95134**

I hereby declare that the above identified small business concern qualifies as a small business concern, as provided in 37 CFR 1.9(d), for purpose of paying reduced fees to the United States Patent and Trademark Office, in that the number of employees of the concern, including those of its affiliates, does not exceed 500 persons.

I hereby declare that the rights under contract or law have been conveyed to, and remain with, the concern identified above with regard to the invention described in

- the specification filed herewith, with title and inventor(s) as listed above.
- the application identified above.
- the patent identified above.

If the rights held by the above identified concern are not exclusive, each individual, concern, or organization having rights in the invention is listed below and no rights to the invention are held by any person, other than the inventor, who would not qualify as an independent inventor under 37 CFR 1.9(c), it that person made the invention, or by any concern which would not qualify as a small business concern under 37 CFR 1.9(d), or a nonprofit organization under 37 CFR 1.9(e).

Each person, concern or organization having any rights in the invention in addition to the above identified concern is listed below:

Name _____
Address _____
 Individual Small Business Nonprofit Organization

I acknowledge the duty to file, in this application or patent, notification of any change in status resulting in loss of entitlement to small entity status prior to paying, or at the time of paying, the earliest of the issue fee

or any maintenance fee due after the date on which status as a small business entity is no longer appropriate. (37 CFR 1.28(b))

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true and further, that these statements were made with knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code, and that such willful false statements may jeopardize the validity of the application, any patent issuing thereon, or any patent to which this verified statement is directed.

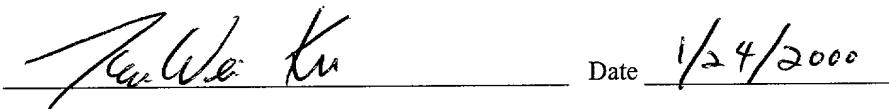
Tsu-Wei Ku
Print Name of Person Signing

Chief Technology Officer
Title of Person Signing

2801 Orchard Parkway, San Jose, CA 95134

Address of Person Signing

SIGNATURE



Date

1/24/2000

*Application
For
United States Non-Provisional Utility Patent*

10 *Title:*

**Method and Apparatus of Relative Datapath Cell Placement with Structure
Bonding and Exponential Field Sampling**

15 *Inventors:*

**Tsu-Wei Ku, residing at 1177 Elmsford Court, Cupertino, CA 95014, a citizen of the
United States**

**Scot A. Woodward, residing at 892 Gallatin Drive #9, Santa Clara, CA 95051, a
citizen of the United States**

**Yung-Hung Wang, residing at 439 Del Medio Avenue, #2, Mountain View, CA
94040, a citizen of Taiwan R.O.C.**

**Duan-Ping Chen, residing at 1303 Peralta Court, San Jose, CA 95120, a citizen of
the United States**

**Wei-Kong Chia, residing at 355 Alicia Way, Los Altos, CA 94022, a citizen of the
United States**

METHOD AND APPARATUS OF RELATIVE DATAPATH CELL PLACEMENT WITH STRUCTURE BONDING AND EXPONENTIAL FIELD SAMPLING

BACKGROUND INFORMATION

5 Field of the Invention

The invention relates to computer-aided design (CAD) of integrated circuits, and particularly to placement of cells in datapath design.

Description of Related Art

Two general categories of placement methods are used, the first one random placement and the second one is structural placement. For random placement, cells are connected by list of nets. The objective of placement is to minimize a cost function that are related to total sum of net length with fixed constraints of placement area, routability, and timing. Several conventional placement algorithms relating to this type of placement method have been disclosed.

In bonding cell instances, a conventional solution places wires at locations between cell instances which produce shortest net lengths. Such approach produces an undesirable random or irregular structure arisen from the placement of uncorrelated cell instances. Another conventional solution places cell instances in a rigid structure by matching pin locations between cell instances. This technique is prohibitively restrictive since adjacent cell instances must be aligned with matching pin locations and hence introduces many dead placement spaces.

In a density map partitioning, conventional solution adopts a uniform grid partitioning. This approach requires substantial computation time in calculating the density forces of a given point on a partition map. Because each square or rectangle in

the grid is uniform in size, the computation time is greater if the grid size of a cell instance is smaller which produces a high number of cells.

Accordingly, it is desirable to place cell instances in an optimized structure bonding. It is further desirable to have a density map partition which is non-uniform in 5 computing a force update vector in reducing the extensive computational complexity time.

SUMMARY OF THE INVENTION

The invention discloses a relative placement of datapath elements by keeping the relativitiy of cell instances in a datapath group without rigid alignments between cells.

10 In one embodiment, to encourage placement of a desirable structure, pseudo cells, pseudo pins, and pseudo nets are selected to be placed at certain locations with respect to real cell instances. The end result produces a cluster of real cell instances that form a desirable structure while minimizing the length of nets. The invention therefore achieves objectives of minimizing the sum of net lengths while encouraging cells instances to form 15 in a desirable structure, such as a column structure, a row structure, or a square structure.

In a further aspect of the invention, a non-uniform partitioning of a density map for calculating a force update vector is disclosed. The partitioning is taking over a region A to compute Riemann sum approximations of a function F over the region A. A force update vector is calculated for a given cell instance within the region A where 20 neighboring cell instances have an exponentially larger grid size as cell instances extend further away from the given cell instance.

Advantageously, the structure bonding reduces the dimension of a chip by minimizing dead placement spaces in a datapath structure for efficient utilization. Moreover, the structure bonding minimizes data skew of datapath. The structure bonding of cell instances with pseudo elements also encourages low timing delays. Additionally, 5 the exponential portioning advantageously reduces significantly on the computational time in the determination of a force update vector on a given cell instances arisen from attractive and repelling forces.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a flow diagram illustrating the general process in the placement of cell 10 instances in accordance with the present invention.

FIGS. 2A-B are pictorial diagrams illustrating relative cell placement of a column structure and a row structure in accordance with the present invention.

FIG. 3 is a pictorial diagram illustrating relative datapath placement in accordance with the present invention.

15 FIG. 4A is a pictorial diagram illustrating cell instances constructed with a pseudo net and pseudo pins in accordance with the present invention; FIG. 4B is a graphical diagram illustrating cell instances with relative positioning in accordance with the present invention.

FIG. 5 is a pictorial diagram of one embodiment in field driven placement with 20 positive and negative charge particles in accordance with the present invention.

FIG. 6 is a flow diagram illustrating the process in a structure bonding in accordance with the present invention.

FIG. 7 is a flow diagram illustrating the process of a Riemann sum approximation of the force update vector in accordance with the present invention.

FIG. 8 is a pictorial diagram illustrating one embodiment of an exponential field sampling partition in accordance with the present invention.

5 DETAILED DESCRIPTION OF PREFERRED EMBODIMENT

FIG. 1 is a flow diagram of illustrating process steps by a placement engine **10** for the placement of cell instances in a structure bonding. Placement engine **10** processes placement steps so that each step generates a new placement based on an existing placement. Each placement step modifies an existing placement by combining attraction actions between connected cell instances with repulsive actions between cell instances. Placement engine **10** receives **11** an input netlist and identifies **12** cell structures from the netlist. Cell structures are extracted **13** from the netlist by placement engine **10** to perform **14** bus analysis. Placement engine **10** then extracts **15** general regularity structures and creates **16** bonding between the extracted cells. Subsequently, placement engine **10** performs **17** of field-driven placement of cell instances on a density partition map.

FIG. 2A is a pictorial diagram illustrating one fundamental relative structure in datapath in a column structure **20**. Column structure **20** embodies a vertical sequence of cell instances being fixed while horizontal distance between two cell instances is not fixed. Preferably in a column structure **20**, cell instances are within the vicinity with one another. For example, in an 8-bit datapath, the binary “0” bit cell is placed at the bottom. The binary “1” cell is then placed above the binary “0” bit cell. This vertical sequence

continues with the binary “2” bit cell being placed on top of the binary “1” cell until all eight cells are stacked in a chronological sequence from the lowest binary value to the highest binary value. Although the 8-bit vertical sequence in the datapath is fixed, cells that are horizontally adjacent to the 8-bit datapath have the flexibility to be placed in the vicinity, without being strictly aligned, to the eight vertical stacked cells. FIG. 2B is a pictorial diagram illustrating another fundamental relative structure in datapath in a row structure **25**. Row structure **25** embodies horizontal sequence of cell instances being fixed while vertical distance between two cell instances is not fixed. A further fundamental relative datapath structure is an array structure where both the horizontal sequence and vertical sequence are fixed. One skilled in the art should recognize that other type of structures or modifications can be practiced without departing from the scope of the invention.

Some terminologies used in describing the invention are defined as follows. A cell instance is an object which inhibits, typically a rectangular subset of a placement area. A net is an object which connects some set of cell instances. A pin is an object which represents the relative location in a cell instance’s region where a net connection occurs. A placement is a mapping from a set of cell instances to a set of positions in a placement area. These positions can represent, for example, the centers or lower left corners of regions that cell instances inhibit. A placement area represents some subset of a two-dimensional real space, which is typically a bounded rectangle or rectilinear region. A total wire length in a placement is the value of some function that approximates the cost associated with the physical wire that connects to physical objects in which cell instances represent. The write length function is chosen to minimize the sum of half-

perimeters of net bounding boxes. For each net, a minimal rectangle is found with horizontal and vertical edges that contain all pins associated that net, and add the width and height of that rectangle to a total wire length function. The utilization μ of a placement is the sum of areas of cell instances, divided by an area of the placement area.

5 The utilization is typically between zero and one, and is independent of the positions of cell instances. A value of zero means that there are no cell instances, and that the placement area is empty. A value of one means that there is exactly enough room for all cell instances, and there would be no unused space if a cell instance can be placed without overlap.

10 FIG. 3 is a pictorial diagram illustrating a structure bonding 30 of cell instances. Structure placement is to place a particular group of cell instances at certain locations that are relative to one another. Cell instances 31a-g on the left side are concatenated to cell instances 35a-g on the right side in a rectangular structure, even though cell instances 31a-g are not perfectly aligned with cell instances 35a-g. Placement engine 10 places 15 cells instances 31a-g and 35a-g with the aid of pseudo nets thereby producing structure bonding 30. Structure bonding 30 results in a zig-zag structure, the usage of pseudo nets encourage cell instances to form a structure without introducing extra dead placement spaces.

FIG. 4A is a pictorial diagram illustrating cell instances constructed with pseudo 20 net and pseudo pins. The shortest distance between a cell instance 40 and a cell instance 41 is a straight line from the top of cell instance 41 to the bottom of cell instance 40. Therefore, a pseudo pin 42 is placed in bottom of cell instance 40 for connecting with a pseudo pin 43 at the top in cell instance 41. A pseudo net 44 connects between pseudo

pin 42 and pseudo pin 43 together while minimizing the length of wire interconnect.

Although pseudo pins 42, 43 and pseudo 44 do not actually exist, pseudo net 44 and pseudo pins 42, and 43 are treated as real nets and pins for encouraging the placement of cell instances in a structure.

5 A pseudo cell instance c_f 48 is added with a width and height equal to zero, without interfering with a density map. The pseudo cell instance c_f 48 is assigned to a position on the density map to encourage a desirable structure. Between pseudo cell instance, c_f and real cell instances c_1 , c_2 , ... c_k , a set of \bar{d}_1 , ..., \bar{d}_k relative positions is created and satisfies: $\bar{x}_j = \bar{x}_f + \bar{d}_j$ for all j, where the term \bar{x}_f denotes an optional 10 pseudo instance and the term \bar{d}_j denotes a relative position.

An example is shown in FIG. 4B. A set of three cell instances c_1 45, c_2 46, and c_3 47, each represents a square of five units in length and width. The three cell instances c_1 45, c_2 46, and c_3 47, are aligned in a cascaded row. If the lower left corner of c_1 45 is at a location \bar{x}_1 , then cell instance c_2 46 has a lower left corner at $\bar{x}_2 = \bar{x}_1 + (5,0)$, and cell 15 instance c_3 47 has a lower left corner $\bar{x}_3 = \bar{x}_1 + (10,0)$. In this example, pseudo cell instance c_f 48 is created for sharing the location \bar{x}_1 with the lower left corner of c_1 . Induce relative position are computes as $\bar{d}_1 = (0,0)$ between c_f 48 and c_1 45, $\bar{d}_2 = (0,5)$ between c_f 48 and c_2 46, and $\bar{d}_3 = (0,10)$ between c_f 48 and c_3 47.

Upon determining the relative positions between the new pseudo cell instance and 20 real cell instances, new pseudo nets are created to bond the new pseudo cell instance and real cell instances. Although several techniques are used to construct pseudo nets, the

preferred technique is to select a connection which minimizes wire lengths. For example, a simple net is used to construct cell instances in a structure. If, however, large nets are given lower weights than small nets, then the bonding effect gained by the net may be diminished. In such scenario, two-pin nets may be created among cell instances. The 5 pins of nets are placed at locations that encourages a desire relative placement of the structure.

In assigning relative locations of pins to their corresponding cell instances, pins associated with the pseudo or fake cell instance c_f 48 are placed at a relative location of (0, 0) from the position of c_f 48. Subsequently, a pin that is associated with a real 10 instance c_j may be assigned to a relative position of \bar{d}_j from c_j . If a relative location of (0, 0) is not assigned to the pseudo cell instance's pin, then the difference of these relative locations is \bar{d}_j , i.e. $\bar{p}_f - \bar{p}_j = \bar{d}_j$ where \bar{p}_f represents the location of c_f 's pin relative to \bar{x}_f , while \bar{p}_j represents the location of c_j 's pin relative to \bar{x}_j .

Continuing with the example above, three fake nets n_1 , n_2 , and n_3 which connect 15 c_f to , respectively, c_1 45, c_2 46, and c_3 47. If a pin is placed for each net at a location of (0, 0) relative to the pseudo cell instance, a pine for n_1 is place at $\bar{d}_1 = (0, 0)$ relative to c_1 , a pine for n_2 is place at $\bar{d}_2 = (0, -5)$ relative to c_2 , and a pine for n_3 is place at $\bar{d}_3 = (0, -10)$ relative to c_3 .

Preferably, placement engine 10 attempts to make pseudo nets as short as possible 20 between cell instances. The shortest possible net is one with a length of zero, in which all of the net's pins are at the same position \bar{z} . For example, if c_j is a real member of such

a net, and if pins are placed at \bar{p}_j relative to \bar{c}_j and \bar{p}_f relative to \bar{c}_f then a desirable result is achieved $\bar{x}_j = \bar{x}_f + \bar{p}_f - \bar{p}_j = \bar{x}_f + \bar{d}_j$, where $\bar{z} = \bar{x}_j + \bar{p}_j = \bar{x}_f + \bar{p}_f$. The optimal structure is to have all pseudo nets have zero length so that cell instances in a structure have exactly the desirable relative positions.

5 Optionally, rather than constructing pseudo cell instances, pin offset vectors can be assigned as if a pseudo instance existed. Each cell instance in the structure is connected to surrounding cell instances by a path of pseudo nets.

Members of a structure may be scattered if real nets connecting to cell instances overpowers pseudo nets. In such scenario, placement engine 10 increases the weight of
10 pseudo nets to reduce the effect of undesirable structure. In one embodiment, the weights of pseudo net are set equal to the weights of real nets to maintain structure cohesion.

FIG. 5 is a flow diagram illustrating the process of structure bonding 50 in accordance with the present invention. A pseudo instance is created 51 and selects 52 at a relative location to encourage structural placement. Placement engine 10 finds 53 offsets between pseudo instance and real instances, represented by the symbol \bar{d}_j . Subsequently, placement engine 10 selects offset between pseudo instance and the associated pseudo pin, \bar{p}_f . After pin offsets are selected 55 for real instances and their associated pseudo pins, \bar{p}_j , placement engine 10 creates a pseudo net connecting pseudo pins.
15

20 FIG. 6 is a pictorial diagram of one embodiment in a field driven placement 60 with positive and negative charge particles. Particles 61 and 62 are tied together with a spring 63. Similarly, particles 64 and 65 are connected with a spring 66. Also, a spring

69 ties particles 67 and 68. Placement engine 10 balances different objectives through a sequence of placement steps. Each placement step modifies the placement by combining attractive actions between connected cells with repulsive actions between cell instances. 5 Attractive actions serve to optimize objective functions while repulsive actions serve to reduce overlap among cell instances. The name “field driven placement” encompasses positively charged particles placed in a negatively charged field. The particles are tie together with springs. The objective is to minimize the following function:

$$\min \sum_{ij} w_{ij} r^2 + \frac{\epsilon}{2} \int E^2 d\tau \quad \text{Eq. (1)}$$

The term r^2 in Equation 1 represents the energy of a spring, which can also be 10 written as $(x_i - x_j)^2 + (y_i - y_j)^2$. The additional term of $\frac{\epsilon}{2} \int E^2 d\tau$ represents the energy of the field. To minimize Equation 1, a vector differentiation is taking on the expression on all movable objects location. The derived Equation 1 becomes: $A\bar{x} + \bar{b} + \bar{F}_x = 0$, and $A\bar{y} + \bar{b} + \bar{F}_y = 0$, where $(A\bar{x} + b, A\bar{y} + b)$ represents the spring force and (F_x, F_y) represents the electrical force.

15 The derived Equation 1 can be interpreted in which the sum of an overall force is equivalent to zero at equilibrium stage. The electrical force is then expressed as follows:

$\bar{F}_x(x, y) = \bar{F}_{xa}(x, y) + \bar{F}_{xr}(x, y)$ and $\bar{F}_y(x, y) = \bar{F}_{ya}(x, y) + \bar{F}_{yr}(x, y)$. The term $F_{xa}(x, y)$ represents the force due to attractive force of the positive plane, while \bar{F}_{xr} represents the 20 repelling force from other cell instances, and thus $\bar{F}_a(x, y) = (\bar{F}_{xa}(x, y), \bar{F}_{ya}(x, y), \bar{F}_r(x, y) = \{\bar{F}_{xr}(x, y), \bar{F}_{yr}(x, y)\})$. Cell instances move in a

plane. According to Gauss's Law: $\oint EdS = q \Rightarrow E = \frac{q\hat{r}}{2\pi r}$. The attraction force F_a

can be calculated as follows: $F_a(x, y) \approx \frac{Ak}{2\pi} \int_i^b \int_j^t \frac{\hat{r}}{r} dxdy$, where symbols l and w represent

the left and right boundaries in a placeable area, and b and t represents the top and bottom
boundaries in the placeable area. Similarly, the repulsive force acting upon a cell is

5 computed as the summation of the repulsive force from other cells: $F_r(x, y) \approx$

$$-\sum_j \frac{Ak}{2\pi} \int_{x_j - \frac{w}{2}}^{x_j + \frac{w}{2}} \int_{y_j - \frac{h}{2}}^{y_j + \frac{h}{2}} \frac{\hat{r}}{r} dxdy.$$

An force update vector is an accumulation of repulsive actions upon a given cell instance by other cell instances. These actions are determined such that cell instances are pushing away from high density regions while pulling towards low density regions. The
10 constant value of k affects the strength of the force update vector. Low values of k reduce the tendency of placement engine 10 to favor overlap reduction over the optimization of objective functions, with a tradeoff against runtime in computing a greater number of force update vectors. High values of k presents a converse issue, in which a poor solution is achieved quickly.

15 Another way to characterize the placement step and the computation of an force update vector is described as follows. An initial placement step is to update a force vector \bar{f}_c for each cell instance c, which is used as an offset vector when finding that instance's location. The force vector \bar{f}_c , which is initially set to zero, is additively

updated during each placement step by repulsive actions from other cell instances, as well as attractive actions from the placement area.

The subsequent placement step comprises an iterative process of a small number of intermediate updates of placement. Each intermediate update include two phases, one which determines virtual positions of the nets and one which determines the positions of the cell instances. The first phase of the intermediate update finds a virtual positions \bar{x}_n for each net n which connects a set of cell instances $\{c_1, \dots, c_j\}$. The position is obtained by computing a weighted average of instances' positions:

$$\bar{x}_n = \frac{\sum_{i=1}^j w_i \cdot \bar{x}_{c_i}}{\sum_{i=1}^j w_i}$$

where $\{w_1, \dots, w_j\}$ are non-negative cell instance weights. A constant weight of 1 is typically used.

During the second phase, an intermediate position \bar{x}_c is found for each cell instance c . This position is a shifted weighted average of the virtual positions of the nets $\{n_1, \dots, n_k\}$ which are connected to c . The weighted average is shifted by the force vector which is computed at the outset of placement steps:

$$\bar{x}_c = \frac{\sum_{i=1}^k w_i \cdot \bar{x}_{n_i}}{\sum_{i=1}^k w_i} + \bar{f}_c$$

where $\{w_1, \dots, w_k\}$ are nonnegative net weights. Typically, the weight w_i is some $(\frac{2}{p})^k$

where p is the pin count of the net n_i (i.e. the number of cell instances which are connected to n_i), and k is some non-negative integer. By starting with $k = 0$ and

gradually increasing to $k = 2$, the total overall wire length is minimized when making transition from a randomly chosen placement to another placement.

Optionally, each net is associated with a force vector which is used as an offset for that net's virtual position, and cell instance positions receive no such offset. This
5 method is referred to as dual placement.

At each placement step, the force vector associated with each cell is updated. The update involves a multiplication by a force trimming scalar first, followed by an addition with a force update vector. The force trimming scalar is typically some real number between zero and one and determines the balance between competing goals of
10 discouraging overlap and optimizing objective functions. Values that are close to 0 encourage shortening nets, while values that are close to 1 encourage cell instances to spread out. The force update vector represents an accumulation of repulsive actions upon a given cell instance by other cell instances. Repulsive actions are determined so that cell instances are pushed away from high density regions and pulled towards low density
15 regions. The force update vector is represented as follows:

$$\bar{e}(\bar{x}) = \frac{k}{2\pi} \iint D(\bar{x}') \cdot \frac{\bar{x} - \bar{x}'}{\|\bar{x} - \bar{x}'\|} d\bar{x}'$$

where k represents some constant, and D represents a density function. The integral is taking over an entire placement area, or all of two-dimensional space as long as the density function is uniformly zero outside the placement area. The force update
20 vector \bar{e}_c associated with a cell instance c is then given by $\bar{e}(\bar{x}_c)$, where \bar{x}_c is the position of c . The constant value k affects the strength of the force update vector. A low

value of k reduces the tendency of placement engine 10 to favor overlap reduction over optimizing objective functions at the expense of runtime. A greater number of force update vectors is computed. A high value of k have the opposite concerns, which achieves poor solutions quickly. In one embodiment, a value of $\frac{3}{1000 \cdot \mu}$ is used where μ

5 represents the utilization of the placement.

A density function D maps a given position \bar{x}' in the placement area to a value which is given by the number of cell instances which cover the position \bar{x}' , less the utilization of the placement: $D(\bar{x}') = \sum_c \chi_c(\bar{x}') - \mu$, where χ_c represents the characteristic function associated with the region that instance c inhabits, i.e. $\chi_c(\bar{x}')$ equals one if the

10 position \bar{x}' is covered by c, and otherwise equals to zero. The subtraction of the utilization from the sum of the characteristic functions serves to provide an attractive action upon cell instances towards regions with low density.

The integral in the force update vector is computed through a Riemann sum approximation. FIG. 7 is a flow diagram of this process 70. The placement engine 10 finds 71 the positions of the sample regions and selects 72 a region. Within that region, the placement engine 10 selects 73 a given point to serve as a representative for computing the part of the force update vector which represents attractive and repulsive actions in that region. The placement engine 10 adds a scaled effect of that point to the overall effect. If there are more regions, the process 70 returns to the operation 72.

15 Otherwise, the placement engine 10 completes the process 70 in a step 76.

To approximate the integral of a function F over some region A , the region A may be partitioned into a collection of pairwise disjoint subregions A_1, \dots, A_k , and then representative points a_1, \dots, a_k may be chosen so that $a_j \in A_j$ for all $j = 1, \dots, k$. If $\|A_j\|$ represents the area of region A_j , then the corresponding Riemann sum approximation

5
$$\int_A F \text{ is: } \sum_{j=1}^k F(a_j) \cdot \|A_j\|.$$

Riemann sums are typically computed by dividing a region A into a rectangular grid, where subregions occupy an equal area. Representative points can be chosen to be corners, centers, random locations, or places where the function to be integrated is minimized or maximized. To make the computation of the force update vector more efficient, the placement area is divided into such a uniform grid, and store a density map that indicates, for each uniform region in the grid, the cell density within that region. This partition is called the density map partition.

There are some drawbacks to using a uniform grid of regions in the computation of the force update vector. Selecting a small grid size promotes accuracy in Riemann sum approximation, but with a severe cost in computation time. Conversely, selecting a large grid size reduces the computational time of force update vectors, but allows for a large amount of irregularity in the cell density within the regions. For example, if 10,000 cell instances are to be placed in an area that is partitioned into two subregions, with 5,000 cell instances for each subregion, then the placement may be categorized as uniform if 20 each of the regions has its 5,000 cell instances occupying the representative position chosen for that region.

To avoid the difficulties arising from the use of a uniform partition in the Riemann sum approximation, a non-uniform partition is chosen whose members are made up of groups of smaller density map grid elements. To this end, a convenient labeling system is established for the density map partition elements base on rows and columns. Given a model rectangle E , the two-dimensional plane is tiled with a grid of copies $\{E_{ij} \mid i,j \in \mathbf{Z}\}$ of E , with the following order: $i_0 < i_1$ if and only if $E_{i_0,j}$ is left of $E_{i_1,j}$ for all j ; and $j_0 < j_1$ if and only if E_{i,j_0} is below of E_{i,j_1} for all i . Without loss of generality, the $E_{0,0}$ region may be placed at any location on the placement area, such as at the lower left corner.

The density map partition induces a map $\delta: \mathbf{Z}^2 \rightarrow \mathfrak{R}$ such that $\delta(i, j)$ is equal to the cell instance density in $E_{i,j}$. For example, if $E_{4,7}$ contains 3 area units of cell instance c_1 , and 2.4 area units of cell instance c_2 , and no other cell instances intersect $E_{4,7}$, then $\delta(4, 7) = 3 + 2.4 = 5.4$. This map, less the utilization μ , is used as the density function D in the Riemann sum calculation of the force update vector.

In describing a non-uniform partition used in Riemann sum approximation, two families of intermediate partitions are described to induce a final partition by way of intersections of intermediate partition members. The intermediate partitions divide a plane into either horizontal or vertical stripes, with a central thin strip that is surrounded by increasingly thick stripes.

Given a density map partition $\{E_{ij} \mid i,j \in \mathbf{Z}\}$, a base index $i_0 \in \mathbf{Z}$, and $m \in \mathbf{Z}$, $V_{i_0}(m)$ can be defined as, the m^{th} vertical stripe about i_0 , as follows:

$$V_{i_0}(0) = \cup \{E_{i_0, j_0} \mid j \in Z\}$$

$$V_{i_0}(m) = \cup \{E_{i_0+k, j} \mid j \in Z, \text{ and } 2^{m-1} \leq k < 2^m\} \text{ for } m > 0$$

$$V_{i_0}(-m) = \cup \{E_{i_0-k, j} \mid j \in Z, \text{ and } 2^{m-1} \leq k < 2^m\} \text{ for } m > 0$$

Similarly given a base index $j_0 \in Z$, and $n \in Z$, $H_{j_0}(n)$ can be defined, the n^{th}

5 horizontal stripe about j_0 , as follows:

$$H_{j_0}(0) = \cup \{E_{i, j_0} \mid i \in Z\}$$

$$H_{j_0}(n) = \cup \{E_{i, j_0+k} \mid i \in Z, \text{ and } 2^{n-1} \leq k < 2^n\} \text{ for } n > 0$$

$$H_{j_0}(-n) = \cup \{E_{i, j_0-k} \mid i \in Z, \text{ and } 2^{n-1} \leq k < 2^n\} \text{ for } n > 0$$

To verify the families of vertical stripes and horizontal strips: let $i_0 \in Z$. The

10 collection $\{V_{i_0}(m) \mid m \in Z\}$ is a partition of a plane. For $m > 0$, the width of $V_{i_0}(m+1)$ is twice that of $V_{i_0}(m)$, and 2^m times that of $V_{i_0}(0)$. For $m < 0$, the width of $V_{i_0}(m-1)$ is twice that of $V_{i_0}(m)$, and 2^{-m} times that of $V_{i_0}(0)$.

In the Riemann sum calculation of a force update vector, the exponential partition is defined as follows. Given a base index $(i_0, j_0) \in Z^2$, and integers m and n ,
15 $A_{i_0, j_0}(m, n)$ can be defined, the $(m, n)^{\text{th}}$ exponential partition member, as
 $V_{i_0}(m) \cap H_{j_0}(n)$. If there is no uncertainty about the values of i_0 and j_0 , then the abbreviation of $A(m, n)$ is used.

The exponential partition members are so named because their areas increase exponentially as the indices m and n move away from zero. Let $(i_0, j_0) \in Z^2$. Then,

$A_{i_0, j_0}(0, 0) = E_{i_0, j_0}$. The collection $\{A(m, n) | m, n \in Z\}$ is a partition of the plane. For $m > 0$ and $n > 0$, the area of $A(m + 1, n + 1)$ is 2^{m+n} times that of $A(0, 0)$. FIG. 8 is a pictorial diagram of one embodiment in implementing an exponential field sampling partition.

The force update vector of a given cell instance c is then computed as described

5 below. Let E_{i_0, j_0} represent the density map partition that contains the center \bar{x} of c .

Representative locations are selected $\{\bar{x}_{m, n} | m, n \in Z\}$ so that $\bar{x}_{m, n} \in A_{i_0, j_0}(m, n)$ for all

m and n in Z . For each m and n in Z , the integers i_m and j_n are found such that $\bar{x}_{m, n}$

$\in E_{i_m, j_n}$. A term $D(m, n)$ is defined to be zero for m and n in Z if E_{i_m, j_n} does not

intersect the placement area; otherwise, $D(m, n)$ is define as $\delta(i_m, j_n) - \mu$, where

10 δ represents the density map and μ represents the utilization of the placement. The resulting Riemann sum approximation is as follows:

$$\iint D(\bar{x}') \cdot \frac{\bar{x} - \bar{x}'}{\|\bar{x} - \bar{x}'\|} d\bar{x}' \approx \sum_m \sum_n D(m, n) \cdot \frac{\bar{x} - \bar{x}_{m, n}}{\|\bar{x} - \bar{x}_{m, n}\|} \cdot \|A(m, n)\|$$

Preferably, representative locations are selected randomly to avoid undesirable clustering of the cell instances.

15 FIG. 9 is a pictorial diagram of an example of one embodiment in implementing an exponential field sampling partition. A given cell instance, denoted by the symbol \diamond , is selected, along with its region in the density map which will serve as a central region. Each of the eight regions immediately surrounding the point \diamond is a 1x1 region unit. Moving outward, the size of a region increases exponentially, from 2 units, to 4 units, and 20 more. The farther away a square or rectangle is from the point \diamond , the larger the sampling

area. Because a single sampling point is chosen for each region, the larger (and thus farther) regions take less processing time per unit of area than the smaller (and closer) regions do. The choice of a single sampling point for a larger region, for example a 4x4 region, also reduces the accuracy of the estimation of that region's effect in the force update vector, but this lack of accuracy is offset by the fact that such regions are far from 5 and, because of the inverse-distance-squared term in the force update vector, have a smaller effect per unit of area on ~~the~~ than a closer region, such as a 2x1 region.

The above embodiments are only illustrative of the principles of this invention and are not intended to limit the invention to the particular embodiments described. For 10 example, variations of pseudo elements may be used to encourage cell instances to be placed in a structure bonding. Furthermore, the exponential partitioning of a density map may have various assigned ratios that are not necessarily powers-of-two. Accordingly, various modifications, adaptations, and combinations of various features of the described embodiments can be practiced without departing from the scope of the invention as set 15 forth in the appended claims.

WE CLAIM:

1 1. A datapath structure, comprising:
2 one or more cell instances, each cell instance having a pin;
3 one or more pseudo cell instances, each pseudo cell instance having a pseudo pin,
4 each pseudo cell instance in the one or more pseudo cell instances being placed at a
5 location relative to the one or more cell instances in encouraging a predetermined
6 structure; and

7 one or more pseudo nets, a first pseudo net connecting between a pin of a first cell
8 instance in the one or more cell instances and a pin in a pin in a first pseudo cell instance
9 in the one or more pseudo cell instances.

1 2. The structure of Claim 1 further comprising a first relative position between
2 the first cell instance and the first pseudo cell instance.

1 3. The structure of Claim 1 wherein the first pseudo cell instance being placed at
2 a location to the first real cell instance thereby producing a zero length in the first pseudo
3 net.

1 4. The structure of Claim 1 wherein the first pseudo cell instance being placed at
2 a location to the first cell instance thereby producing the first pseudo having a value
3 which is greater than a zero length.

1 5. The structure of Claim 1 wherein the predetermined structure comprises a
2 column structure, a row structure, or a square structure.

1 6. A datapath structure, comprising:
2 in a datapath structure, a first cell placed at a first position; and
3 a second cell being placed relatively at a second position to the first position.

1 7. The datapath structure of Claim 6 wherein the second cell being relatively
2 placed such that the first position of the first cell is not strictly aligned to the second
3 position of the second cell.

1 8. A datapath structure of Claim 6 further comprising a pseudo element for
2 aiding in relative placement of the second cell at the second position to the first cell at the
3 first position.

1 9. A datapath structure of Claim 6 wherein the datapath structure comprises a
2 column structure with a fixed vertical sequence for placing the first cell and the second
3 cell.

1 10. A datapath structure of Claim 6 wherein the datapath structure comprises a
2 row structure with a fixed horizontal sequence for placing the first cell and the second
3 cell.

1 11. A datapath structure of Claim 6 wherein the datapath structure comprises an
2 array structure with a fixed vertical sequence and a fixed horizontal sequence.

1 12. A computerized method for encouraging a structure bonding, comprising the
2 steps of:

3 placing a first pseudo cell instance at a location relative to a first cell
4 instance in a plurality of cell instances for encouraging a predetermined structure
5 bonding in the plurality of cell instances; and

6 connecting the pseudo net between the cell instance and the pseudo cell
7 instance.

1 13. The method of Claim 12 further comprising the step of minimizing a wire
2 length in the pseudo net from the placement of the first pseudo cell instance relative to the
3 first cell instance.

1 14. The method of Claim 12 further comprising the step of providing a first offset
2 between the pseudo cell instance and the first cell instance.

1 15. The method of Claim 12 further comprising the step of determining a second
2 offset between the pseudo cell instance and a second cell instance in the plurality of cell
3 instances.

1 16. The method of Claim 12 wherein the predetermined structure comprises a
2 column structure, a row structure, or a square.

1 17. The method of Claim 12 wherein the placing step comprises the step of
2 placement without introducing extra dead placement spaces.

1 18. A density map partition having a region A for computing a force update
2 vector, the region A having a plurality of cell instances with a centering cell at an A(0, 0)
3 location, comprising:

4 a first cell instance density at an A(0, 0) location having a rectangular grid
5 unit; and

6 a plurality of rectangles A(m, n) cell instances coupled to the A(0, 0), the
7 plurality of rectangles A(m, n) cell instances contains multiple number of the
8 rectangular grid unit wherein a farther away A(m, n) cell instance the large the
9 multiple number of the rectangular grid unit.

1 19. The density map partition of Claim 18 wherein the A(m, n) cell instances
2 comprises A(-1, 0), A(-1, 1), A(-1, -1), A(0, 1), A(0, -1), A(1, 0), A(1, 1), A(1, -1) cell
3 instances wherein each having a same rectangular grid unit as A(0, 0).

1 20. The density map partition of Claim 18 wherein the A(m, n) cell instances
2 comprises A(-2, 0), A(-2, 1), A(-2, -1), A(2, 0), A(2, -1), A(2, 1), A(-1, -2), A(0, -2), A(-
3 1, -2), A(-1, 2), A(0, -2), A(1, 2), cell instances wherein each having twice the
4 rectangular grid unit as A(0, 0).

1 21. The density map partition of Claim 18 wherein the force update vector
2 comprises computing attractive and repelling forces affecting the A(0, 0) cell instance.

1 22. A computerized method for generating non-uniform partitioning of cell
2 instances in computing force update vector, comprising the steps of:

3 selecting a reference cell instance in a region A having a plurality of cell
4 instances, the reference cell instance having a grid base unit; and
5 computing a force update vector of the reference cell instance, each of the
6 plurality of cell instances having either a same grid base unit or a multiple time of the
7 grid base unit.

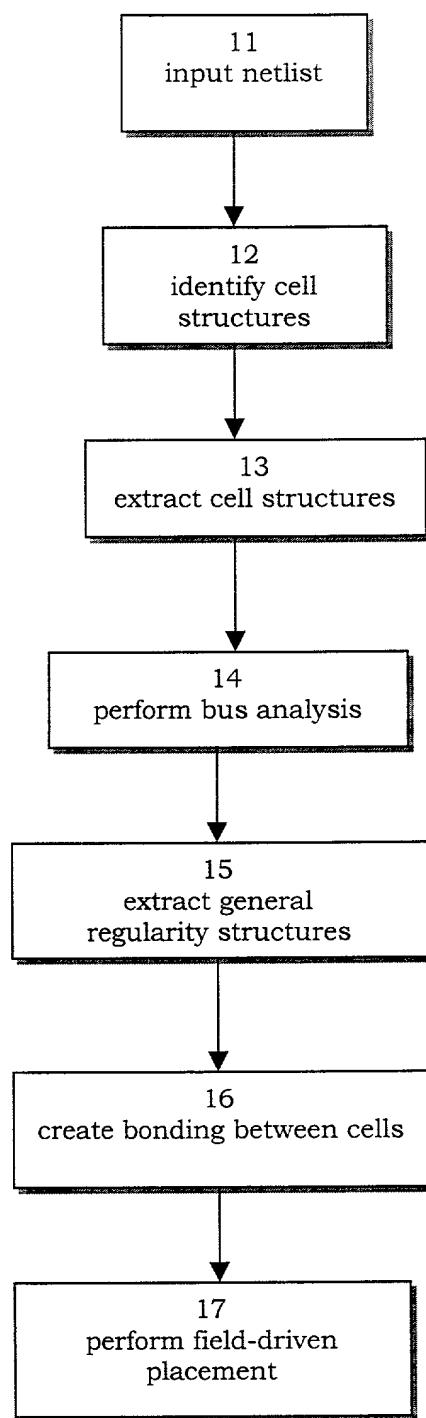
1 23. The method of Claim 22 further comprising the step of computing an
2 attractive force from the reference cell instance in the plurality of cell instances.

1 24. The method of Claim 22 further comprising the step of computing a repulsive
2 force from the reference cell instance in the plurality of cell instances.

ABSTRACT

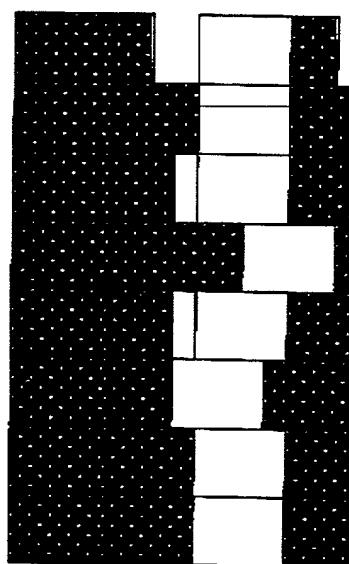
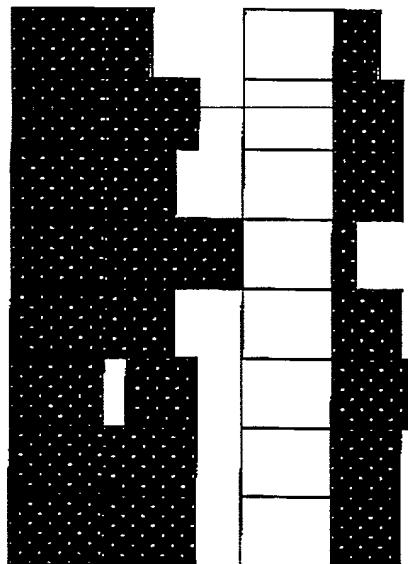
The invention discloses a relative structure placement of datapath of cell instances in a column structure, a row structure, or an array structure. To encourage placement of a desirable structure, pseudo cells, pseudo pins, and pseudo nets are selected to be placed at certain locations with respect to real cell instances. The end result produces a cluster of real cell instances that form a desirable structure while minimizing the length of nets.

The invention further discloses a non-uniform partitioning of a density map for calculating a force update vector. The partitioning is taking over a region A to compute Riemann sum approximations of a function F over the region A. A force update vector is calculated for a given cell instance within the region A where neighboring cell instances have an exponentially larger grid size as cell instances extend further away from the given cell instance.



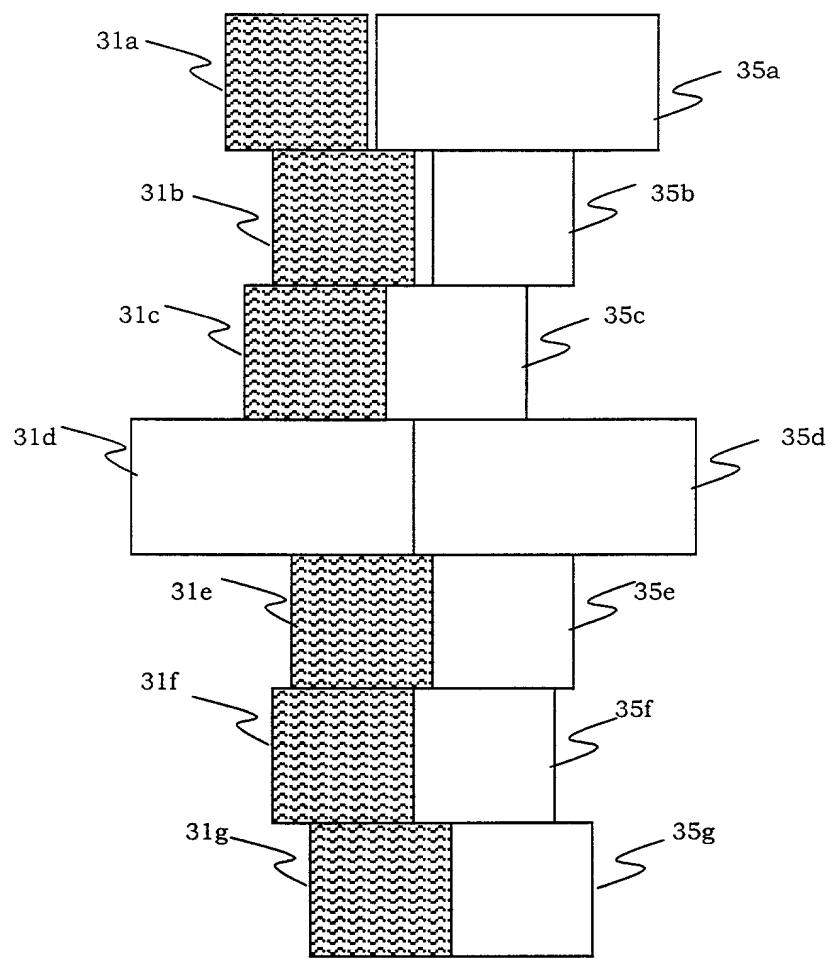
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FIGURE 2A



25

FIGURE 2B



30

FIGURE 3

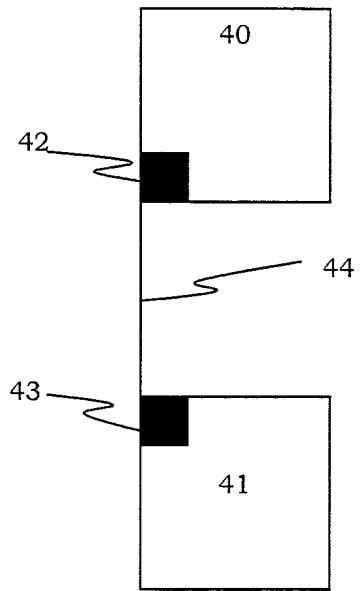


FIGURE 4A

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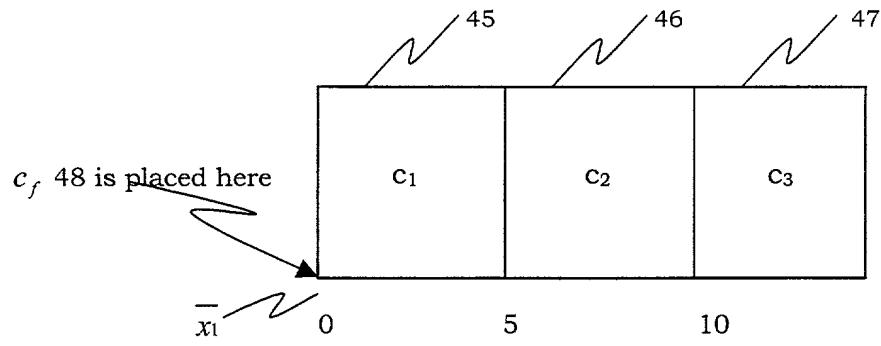


FIGURE 4B

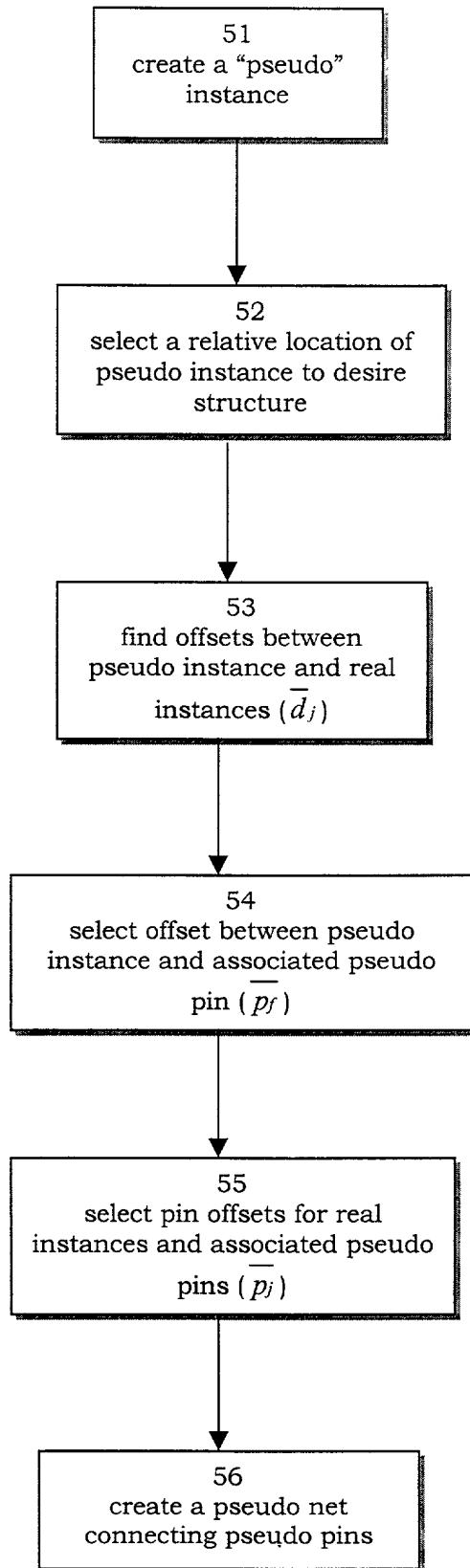
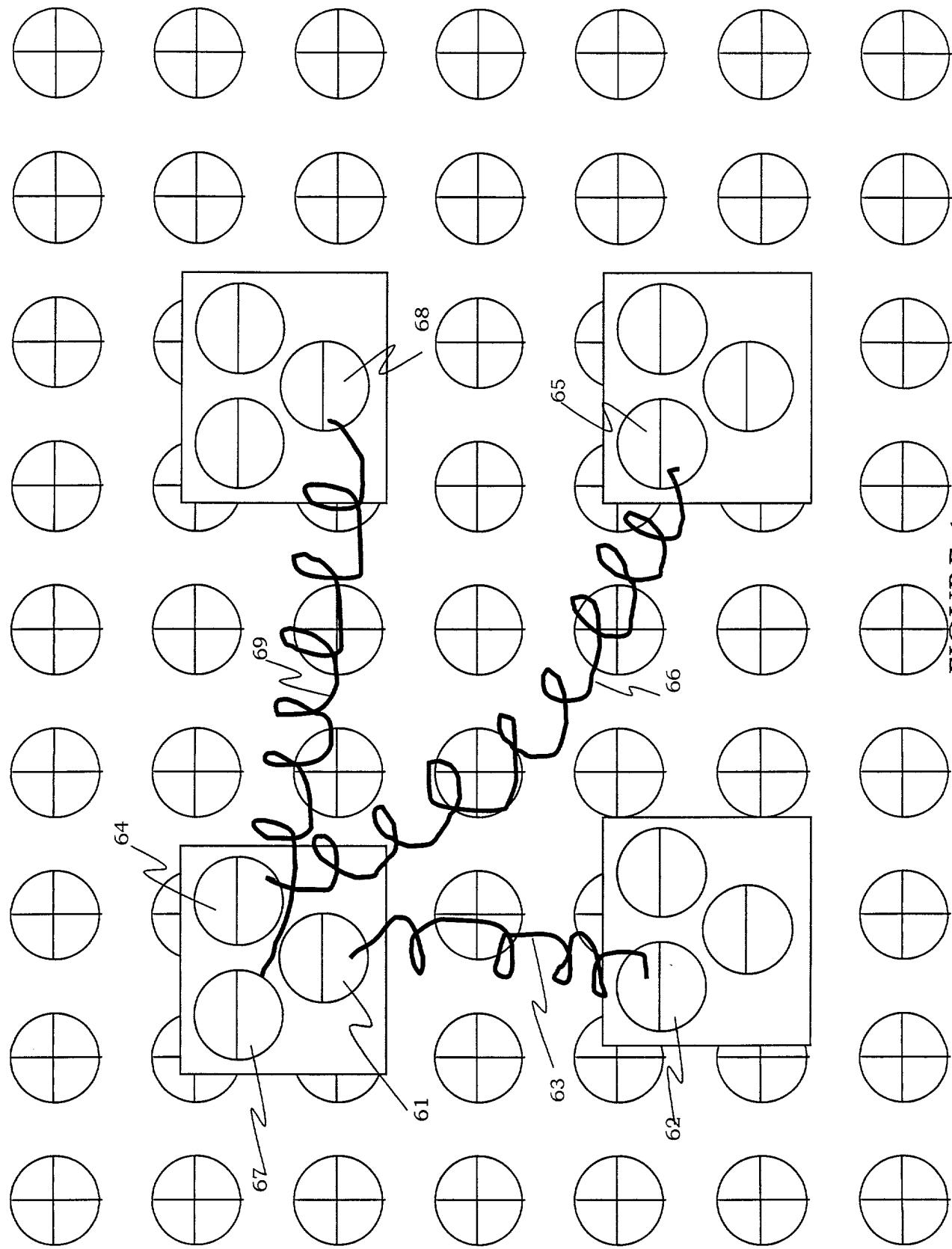
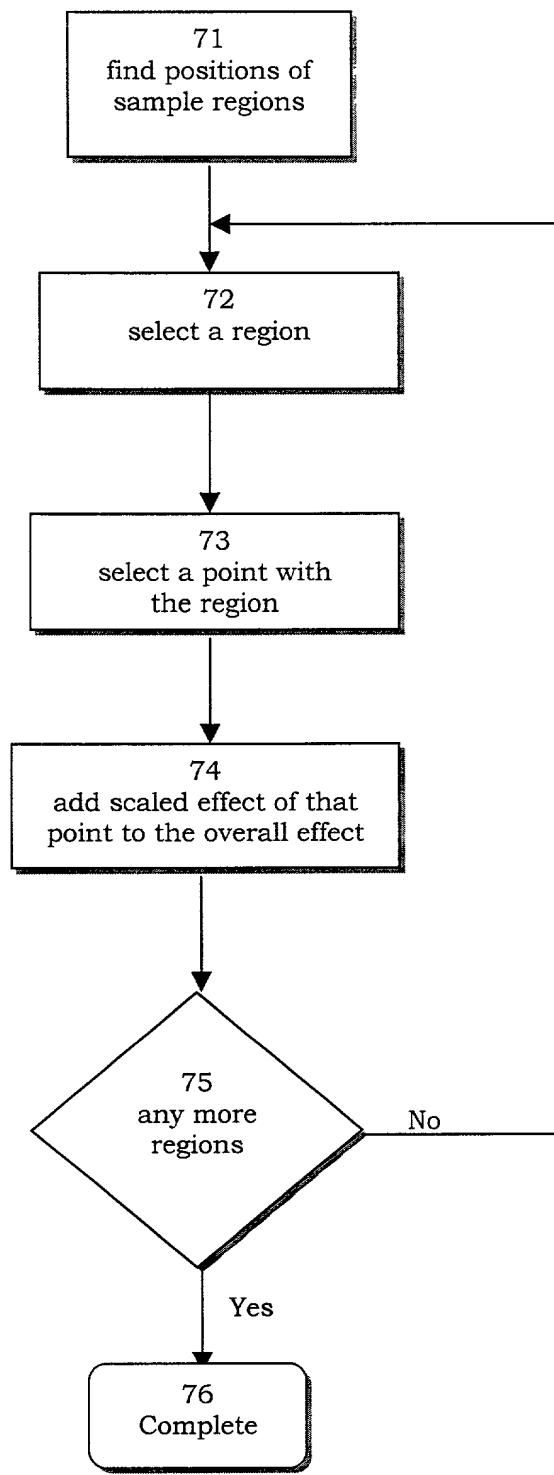


FIGURE 6





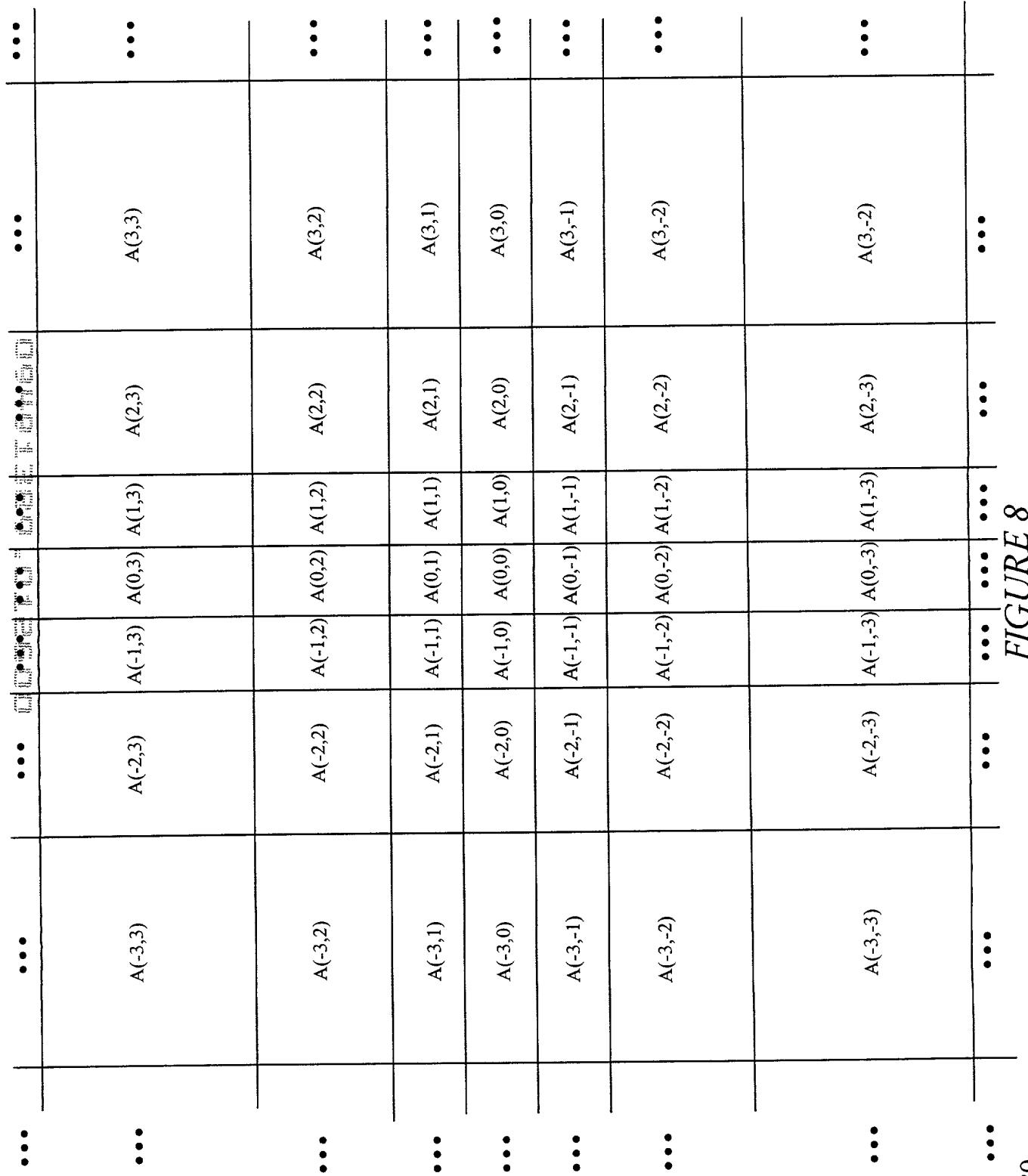


FIGURE 8

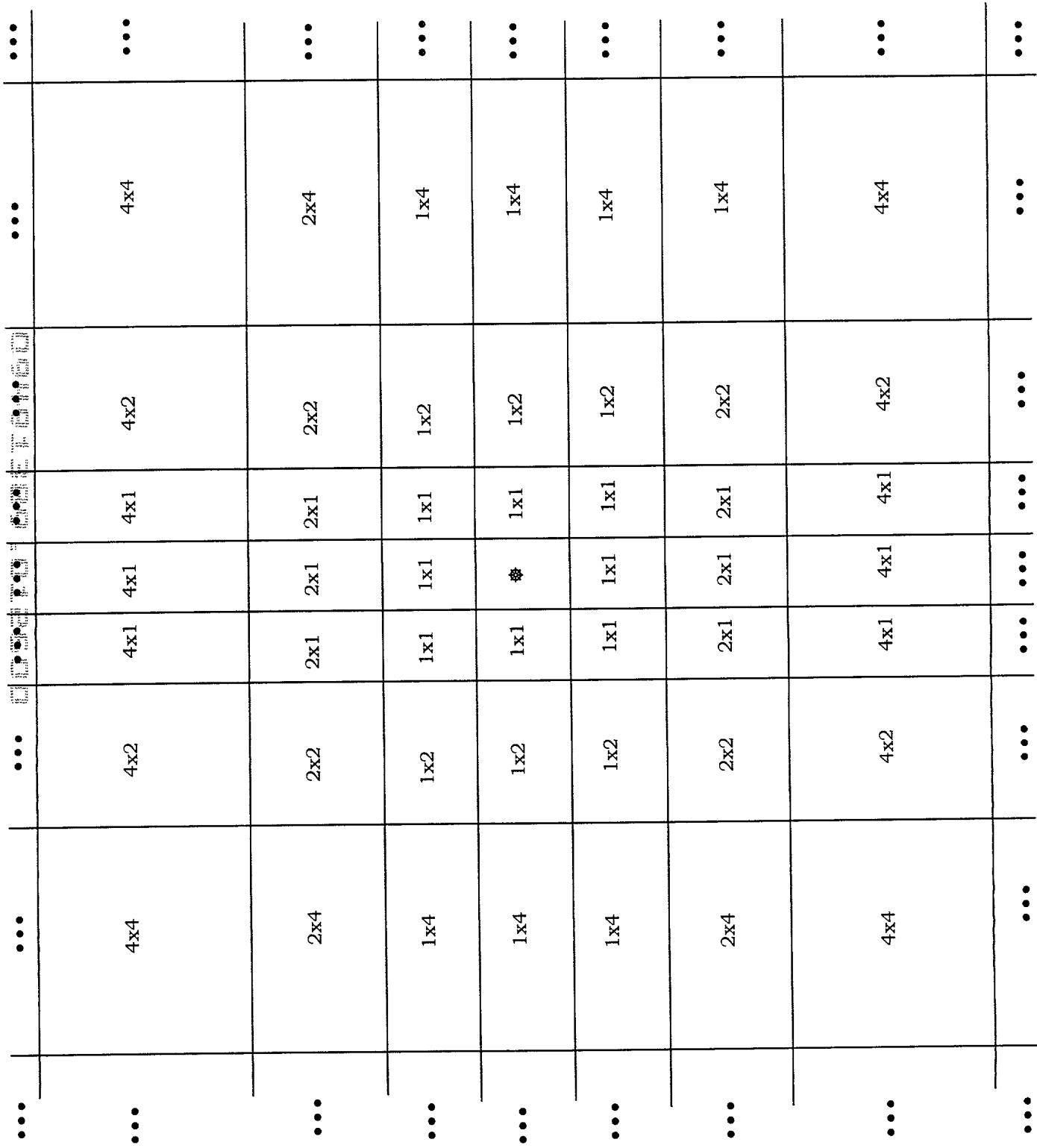


FIGURE 9

DECLARATION AND POWER OF ATTORNEY

As a below named inventor, I hereby declare that:

My residence, postal address and citizenship are as stated below next to my name.

I believe that I am the original, first and sole inventor (if only one name is listed below) or an original first and joint inventor (if multiple names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled:

**Method and Apparatus of Relative Datapath Cell Placement with Structure Bonding and
Exponential Field Sampling**

the specification of which is attached hereto unless the following information is indicated:

— was filed on _____;
as United States Application Number or PCT International Application Number _____;
and was amended on _____ (if applicable).

I hereby state that I have reviewed and understand the contents of the above identified specification, including the claims, as amended by any amendments referred to above.

I acknowledge the duty to disclose information that is material to patentability as defined in 37 CFR 1.56.

I hereby claim foreign priority benefits under 35 U.S.C. 119(a)-(d) or 365(b) of any foreign application(s) for patent or inventor's certificate, or 365(a) of any PCT International application which designated at least one country other than the United States, listed below and have also identified, as so indicated below, any foreign application for patent or inventor's certificate, or PCT International application having a filing date before that of the application on which priority is claimed.

Prior Foreign Application(s)	Priority Claimed
_____	____ Yes ____ No
(Application No.)	(Country)
(Day/Month/Year Filed)	

I hereby claim the benefit under 35 U.S.C. 119(e) of any United States provisional application(s) listed below.

(Application No.)	(Day/Month/Year Filed)
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I hereby claim the benefit under 35 U.S.C. 120 of any United States application(s), or 365(c) of any PCT International application designating the United States, listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States or PCT International application in the manner provided by the first paragraph of 35 U.S.C. 112, I acknowledge the duty to disclose information that is material to patentability as defined in 37 C.F.R. 1.56 that became available between the filing date of the prior application and the national or PCT International filing date of this application.

(Application No.)	(Day/Month/Year Filed)	(Status - patented, pending, abandoned)
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I hereby appoint the following attorney(s) and/or agent(s) to prosecute this application and to transact all business in the United States Patent and Trademark Office connected therewith:

**DENNIS S. FERNANDEZ, REG. NO. 34,160
IRENE H. FERNANDEZ, REG. NO. 34,625
PETER C. SU, REG. NO. 43,939**

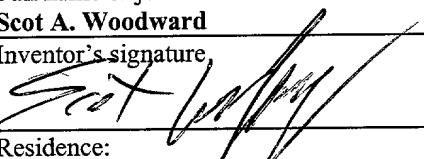
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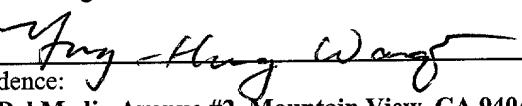
FERNANDEZ & ASSOCIATES LLP
PATENT ATTORNEYS
PO BOX D
MENLO PARK, CA 94026-6204

(650) 529-9237
(650) 529-9315 : FAX
EMAIL: *iploft@iploft.com*

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and beliefs are believed to be true; and further that these statements were made with knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

Full name of sole or first inventor: Tsu-Wei Ku	
Inventor's signature 	Date <i>1/24/2000</i>
Residence: 1177 Elmsford Court, Cupertino, CA 95014	
Citizenship: United States	
Postal Address: SAME AS RESIDENCE	

Full name of joint inventor: Scot A. Woodward	
Inventor's signature 	Date <i>23 Jan 2000 ce</i>
Residence: 892 Gallatin Drive #9, Santa Clara, CA 95051	
Citizenship: United States	
Postal Address: SAME AS RESIDENCE	

Full name of joint inventor: Yung-Hung Wang	
Inventor's signature 	Date <i>Jan 24, 2000</i>
Residence: 439 Del Medio Avenue #2, Mountain View, CA 94040	
Citizenship: Taiwan, R.O.C.	
Postal Address: SAME AS RESIDENCE	

Full name of joint inventor:	
Duan-Ping Chen	
Inventor's signature	Date
	01/24/2000
Residence:	
1303 Peralta Court, San Jose, CA 95120	
Citizenship:	
United States	
Postal Address:	

Full name of joint inventor:	
Wei-Kong Chia	
Inventor's signature	Date
	1/24/2000
Residence:	
355 Alicia Way, Los Altos, CA 94022	
Citizenship:	
United States	
Postal Address:	